

	$z$	$x_1$	$x_2$	$x_3$	$x_4$	RHS
$z$	1	-2	0	-3	0	-12
$x_4$	0	-3	0	-2	1	0
$x_2$	0	1	1	1	0	4

To find the range  $[2, \lambda_2]$  over which this tableau is optimal, we first find  $\bar{\mathbf{b}}$  and  $\bar{\mathbf{b}}'$ :

$$\bar{\mathbf{b}} = \mathbf{B}^{-1}\mathbf{b} = \begin{bmatrix} -2 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 6 \\ 6 \end{bmatrix} = \begin{bmatrix} -6 \\ 6 \end{bmatrix}$$

$$\bar{\mathbf{b}}' = \mathbf{B}^{-1}\mathbf{b}' = \begin{bmatrix} -2 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} -1 \\ 1 \end{bmatrix} = \begin{bmatrix} 3 \\ -1 \end{bmatrix}.$$

Therefore,  $S = \{2\}$  and  $\lambda_2$  is given by

$$\lambda_2 = \frac{\bar{b}_2}{-\bar{b}'_2} = \frac{6}{-(-1)} = 6.$$

For  $\lambda$  in the interval  $[2, 6]$  the optimal objective value and the right-hand-side are given by

$$z(\lambda) = \mathbf{c}_B \bar{\mathbf{b}} + \lambda \mathbf{c}_B \bar{\mathbf{b}}'$$

$$= (0, -3) \begin{pmatrix} -6 \\ 6 \end{pmatrix} + \lambda(0, -3) \begin{pmatrix} 3 \\ -1 \end{pmatrix} = -18 + 3\lambda$$

$$\bar{\mathbf{b}} + \lambda \bar{\mathbf{b}}' = \begin{bmatrix} -6 \\ 6 \end{bmatrix} + \lambda \begin{bmatrix} 3 \\ -1 \end{bmatrix} = \begin{bmatrix} -6 + 3\lambda \\ 6 - \lambda \end{bmatrix}.$$

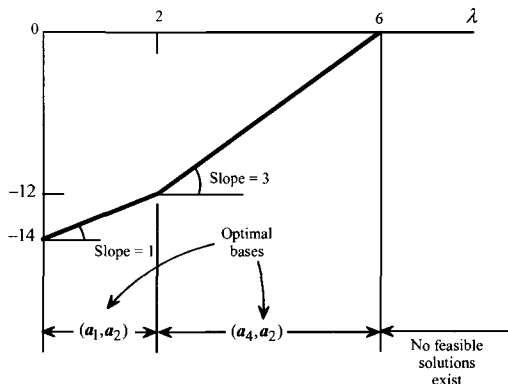
The optimal tableau over the interval  $[2, 6]$  is depicted below:

	$z$	$x_1$	$x_2$	$x_3$	$x_4$	RHS
$z$	1	-2	0	-3	0	$-18 + 3\lambda$
$x_4$	0	-3	0	-2	1	$-6 + 3\lambda$
$x_2$	0	1	1	1	0	$6 - \lambda$

At  $\lambda = 6$ ,  $x_2$  drops to zero. Since all entries in the  $x_2$  row are nonnegative, we stop with the conclusion that no feasible solutions exist for all  $\lambda > 6$ . Figure 6.8 summarizes the optimal bases and the corresponding objective values for  $\lambda \geq 0$ . Note that the optimal objective value as a function of  $\lambda$  is piecewise linear and convex. In Exercise 6.66 we ask the reader to show that this is always true. The breakpoints correspond to the values of  $\lambda$  for which alternative optimal dual solutions exist.

**Comment on Deriving Shadow Prices via a Parametric Analysis**

Observe that parametric analysis can be used to ascertain the structure of the optimal value function  $z^*(b_i)$  (see Equation (6.4)) as a function of  $b_i$ , in the neighborhood of the current value of  $b_i$ , for any  $i \in \{1, \dots, m\}$ . Accordingly, we can then determine the right-hand and left-hand shadow prices with respect to



**Figure 6.8. Optimal objectives and bases as a function of  $\lambda$ .**

$b_i$  as the respective right-hand and left-hand derivatives of  $z^*(b_i)$  at the current value of  $b_i$ , where the former value is taken as infinity in case an increase in  $b_i$  renders the primal problem in Equation (6.1) infeasible.

More specifically, consider determining the right-hand shadow price with respect to  $b_i$ , for some  $i \in \{1, \dots, m\}$ . In this case, the right-hand-side  $\mathbf{b}$  is replaced by  $\mathbf{b} + \lambda \mathbf{b}'$ , where  $\mathbf{b}' = \mathbf{e}_i$ , the  $i$ th unit vector. Accordingly, we can now perform the foregoing parametric analysis until we arrive at a tableau that remains optimal as  $\lambda$  increases from zero (up to some positive level) or else, we detect unboundedness of the dual (infeasibility of the primal) as  $\lambda$  increases from zero. In the former case, the right-hand shadow price is given by  $w_i = (\mathbf{c}_B \mathbf{B}^{-1})_i$  for the corresponding current tableau, and in the latter case, it is infinite in value. In a similar manner, we can compute the left-hand shadow price as the  $w_i$  value corresponding to the tableau that remains optimal as  $\lambda$  increases from the value of zero, where the right-hand-side is now perturbed according to  $\mathbf{b} - \lambda \mathbf{e}_i$ . Exercise 6.70 asks the reader to illustrate this approach.

**EXERCISES**

[6.1] Use the standard form of duality to obtain the dual of the following problem. Also verify the relationships in Table 6.1.

$$\begin{array}{ll}
 \text{Minimize} & \mathbf{c}_1 \mathbf{x}_1 + \mathbf{c}_2 \mathbf{x}_2 + \mathbf{c}_3 \mathbf{x}_3 \\
 \text{subject to} & \mathbf{A}_{11} \mathbf{x}_1 + \mathbf{A}_{12} \mathbf{x}_2 + \mathbf{A}_{13} \mathbf{x}_3 \geq \mathbf{b}_1 \\
 & \mathbf{A}_{21} \mathbf{x}_1 + \mathbf{A}_{22} \mathbf{x}_2 + \mathbf{A}_{23} \mathbf{x}_3 \leq \mathbf{b}_2 \\
 & \mathbf{A}_{31} \mathbf{x}_1 + \mathbf{A}_{32} \mathbf{x}_2 + \mathbf{A}_{33} \mathbf{x}_3 = \mathbf{b}_3 \\
 & \mathbf{x}_1 \geq \mathbf{0} \\
 & \mathbf{x}_2 \leq \mathbf{0} \\
 & \mathbf{x}_3 \quad \text{unrestricted.}
 \end{array}$$

[6.2] Give the dual of the following problem:

$$\begin{array}{rcll}
 \text{Maximize} & -2x_1 + 3x_2 + 5x_3 & & \\
 \text{subject to} & -2x_1 + x_2 + 3x_3 + x_4 & \geq & 5 \\
 & 2x_1 & + & x_3 & = & 4 \\
 & & -2x_2 + x_3 + x_4 & \leq & 6 \\
 & & & x_1 & \leq & 0 \\
 & & & x_2, & x_3 & \geq & 0 \\
 & & & & & x_4 & \text{unrestricted.}
 \end{array}$$

[6.3] Consider the following problem:

$$\begin{array}{rcl}
 \text{Maximize} & -x_1 + 3x_2 & \\
 \text{subject to} & 2x_1 + 3x_2 \leq 6 \\
 & x_1 - 3x_2 \geq -3 \\
 & x_1, \quad x_2 \geq 0.
 \end{array}$$

- Solve the problem graphically.
- State the dual and solve it graphically. Utilize the theorems of duality to obtain the values of all the primal variables from the optimal dual solution.

[6.4] Solve the following linear program by a graphical method:

$$\begin{array}{rcl}
 \text{Maximize} & 3x_1 + 3x_2 + 21x_3 & \\
 \text{subject to} & 6x_1 + 9x_2 + 25x_3 \leq 15 \\
 & 3x_1 + 2x_2 + 25x_3 \leq 20 \\
 & x_1, \quad x_2, \quad x_3 \geq 0.
 \end{array}$$

(Hint: Utilize the dual problem.)

[6.5] Consider the following problem:

$$\begin{array}{rcl}
 \text{Maximize} & 10x_1 + 24x_2 + 20x_3 + 20x_4 + 25x_5 & \\
 \text{subject to} & x_1 + x_2 + 2x_3 + 3x_4 + 5x_5 \leq 19 \\
 & 2x_1 + 4x_2 + 3x_3 + 2x_4 + x_5 \leq 57 \\
 & x_1, \quad x_2, \quad x_3, \quad x_4, \quad x_5 \geq 0.
 \end{array}$$

- Write the dual problem and verify that  $(w_1, w_2) = (4, 5)$  is a feasible solution.
- Use the information in Part (a) to derive an optimal solution to both the primal and the dual problems.

[6.6] Consider the following problem:

$$\begin{array}{rcl}
 \text{Minimize} & 2x_1 + 15x_2 + 5x_3 + 6x_4 & \\
 \text{subject to} & x_1 + 6x_2 + 3x_3 + x_4 \geq 2 \\
 & -2x_1 + 5x_2 - 4x_3 + 3x_4 \geq -3 \\
 & x_1, \quad x_2, \quad x_3, \quad x_4 \geq 0.
 \end{array}$$

- Give the dual linear problem.
- Solve the dual geometrically.
- Utilize information about the dual linear program and the theorems of duality to solve the primal problem.

[6.7] Consider the following linear programming problem:

$$\begin{array}{rcl}
 \text{Maximize} & 2x_1 + 3x_2 + 5x_3 & \\
 \text{subject to} & x_1 + 2x_2 + 3x_3 \leq 8 \\
 & x_1 - 2x_2 + 2x_3 \leq 6 \\
 & x_1, \quad x_2, \quad x_3 \geq 0.
 \end{array}$$